Optimizing Information Mediators by Selectively Materializing Data

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Information Mediators



Movies

Talk Outline

- Performance speed of application dependent on sources
- Approach to performance optimization by local materialization
- Materialization framework for mediators
- Design of materialization system
- Selecting data to materialize
 - Distribution of user queries
 - Structure of sources
 - Updates
- Admission and replacement
- ♦ The integrated materialization system
- ♦ Experimental results
- Related work, applicability to other mediator systems
- Conclusion and future directions

Performance Issue in Information Mediators

- Speed of the application is heavily dependent on sources
- Query response time is high despite having high quality query plans
- Dominant cost is retrieving data from remote sources
 - May have to retrieve a large number of Web pages
 - Source is structured such that retrieving data is time consuming
 - Source may be slow
- Typical Query: "Find all chinese restaurants in Santa Monica with an excellent food rating"
- Takes several minutes to return an answer

Solution: Materialize Data Locally

- ♦ Materialize data locally
- Materializing all the data is impractical
 - Mediator degenerates into data warehouse
- Significant performance gain can be achieved by materializing small fraction of data
 - Hypotheses that some portions of data queried more frequently
 - Materializing certain portions of data speeds up response time for expensive queries
- Data has to be *selectively* materialized
- Primary Issues
 - How is materialized data represented and used
 - How do we automatically identify what to materialize

Overall Approach: Define Materialized Data as Another Information Source



• Existing mediator infrastructure to address two issues

- Providing semantic description of materialized data contents
- Query planner can reason with contents of materialized data

Selecting Data to Materialize



Materialization System : Architecture



Distribution of User Queries: Extracting Patterns

SELECT name, tel FROM restaurant WHERE cuisine="Chinese"

SELECT name, review, address FROM restaurant WHERE city="Los Angeles"

SELECT name, address FROM restaurant WHERE cuisine="Mexican"

SELECT name, tel, address FROM restaurant WHERE cuisine="Chinese"

SELECT name, review FROM restaurant WHERE cuisine="Italian"

SELECT name, address FROM restaurant WHERE city="Santa Monica"

SELECT name, tel, review



CM Algorithm for Extracting Patterns

Too many classes i.e, new information sources create performance problems for query planner

Compact description of patterns extracted

Analyze each query in query distribution
Create subclasses of interest by analyzing constraints
For each subclass cluster attribute groups
Merge across class coverings
Outputs compact description



Analyze constraints in each query

- ♦ Identify subclasses of information of interest
- Maintain ontology in KR system LOOM
- Record attribute groups queried for each subclass

Clustering Attribute Groups

Santa Monica

(name, address, showtimes) 13
(movieurl, tel) 12
(tel, reviews, name) 5
(name, showtimes) 2
(name, address) 2
(movieurl, tel, reviews) 4
(tel, reviews) 7
(name, showtimes, trailers) 10
(name, showtimes) 8

(name, address, showtimes) 13
(name, showtimes) 8
(name, showtimes, trailers) 10
(name, showtimes) 2
(name, address) 2
(tel, reviews, name) 5
(tel, reviews) 7
(movieurl, tel, reviews) 4
(movieurl, tel) 12

• Cluster by attribute group similarity and hits

◆ 2D clustering - optimal clustering NP complete, approximate

Clustering Attribute Groups

Santa Monica

(name, address, showtimes) 13 (movieurl, tel) 12 (tel, reviews, name) 5 (name, showtimes) 2 (name, address) 2 (movieurl, tel, reviews) 10 (tel, reviews) 7 (name, showtimes, trailers) 10 (name, showtimes) 8 (name, address, showtimes, trailers) 10

(name, address, showtimes) 2

(tel, reviews, name) 6

(movieurl, tel, reviews) 11

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Covering: (chinese, mexican, italian) --> Restaurant
(chinese, {A}) U (mexican, {A}) U (italian, {A}) --> (Restaurant, {A})

Merging Across Coverings



Effectiveness, Complexity



- Measured 'precision' and 'recall' in extracting patterns
- Pattern P in query distribution
 - Precision is % of patterns extracted that is in P
 - Recall is % of P that is in patterns extracted
- ◆ High precision and recall for q=0.2
- Complexity = $O(M^2N^2)$
 - M = number of queries, N = Number of attributes in a class

Source Structure Analysis

Problem: Certain kinds queries are expensive as wrapped Web sources not originally designed for database like querying
 Solution: Prefetch and materialize data to improve response time
 Such data cannot be identified by analyzing user queries



GUI Specification

- ◆ Mediator GUI is typically more restrictive
- Formal specification language
- ◆ Data items that can be retrieved
- ◆ Details of selection conditions that can be specified
- SELECT {name, tel, address, cuisine, review, city, rating, map}
 FROM ent

WHERE [city,1,(LA, NYC, Santa Monica)] {cuisine,1,(chinese,...)}

Query Processing Axioms

Precompiled axioms for query processing

restaurant(name,cuisine,address,tel)= zagats(z.name,z.cuisine,z.address,z.tel)

restaurant(name,cuisine,address,tel,lat,long)= zagats(z.name,z.cuisine,z.address,z.tel) and ent_geocoder(\$z.address,g.lat,g.long)

- Axioms tell what data operations will be performed on what sources
- Can be used to determine data to prefetch
- Cost Estimator: Costs of queries
- Process of Source Structure Analysis
 - Use GUI specification and axioms to identify queries
 - Use cost estimator to determine expensive queries
 - Use axioms and knowledge of type of query to determine data to prefetch

Source Structure Analysis

Example :

GUI specification : selection queries on "cuisine" of restaurant Cost estimator : Expensive query

Query processing axioms:

restaurant(name,cuisine,address,tel)= zagats(z.name,z.cuisine,z.address,z.tel)

Heuristic : Prefetch key (name) and selection attribute (cuisine)

Optimization : selection can now be done locally, thus faster

• Examples of heuristics

- 1. selection query materialize key and selection attribute
- 2. join query materialize join attributes and keys
- 3. ordered join materialize result of ordered join

Updates

- Data materialized can change at original sources
- Strategy
 - Do not materialize very frequently updated data
 - Refresh materialized data at appropriate intervals
- Specifying update characteristics, frequency
- Need not assume that user always absolutely requires the latest data
- Also specify user's requirements for freshness of data



Update Specifications

CLASS	MEMBERSHIP	CHANGE	TIME_PERIOD	TIME
MOVIE_SRC	Α	Y	1 week	week : friday

ATTRIBUTE	CHANGE	TIME_PERIOD	TIME
actors	Ν	-	-

CLASS	TOLERANCE
MOVIE	0

ATTRIBUTE	TOLERANCE
theatre	0
actors	0
review	6 weeks

Maintenance Frequency and Cost



For each attribute:
 maintenance period = max(update period, tolerance)
 = tolerance , if update period unknown

♦ Example

review = max(1 week, 6 weeks) = 6 weekstheatre = max(1 week, 0) = 1 week

Maintenance frequency and Query cost - > Maintenance cost

◆ Do not materialize an attribute if maintenance cost too high

◆ Total maintenance cost must be within limit

Admission and Replacement

- Choose optimal set of classes to materialize under 2 constraints
 - Space and Maintenance Cost
- Same as fractional knapsack in 2 dimensions
 - Greedy algorithm for multi-dimensional fractional knapsack



- w_i (query response time savings), s_i (space), m_i (maintenance cost)
- Store in order of wi/max(si,mi)
- ◆ Guarantees optimal selection, generalizes to n dimensions

Materialization System : Architecture



Optimizer Module

- Gets set of classes of data to materialize
 Retrieves data from original sources
 Populates local database with data
 Makes appropriate changes to domain and source models
- ♦ Recompiles axioms
- Periodically refreshes materialized classes
- Periodically reevaluates materialized data classes

Experimental Results

Query Set	Response Time (No optimization)	Response Time (with Optimizer)	%improvement
Q1	38115	9301	75 %
Q2	44186	3775	91 %

(a) Source structure analysis (Countries)

Query Set	Response Time	Response Time	Response Time	%improvement	%improvement
	(No optimization)	(with Optimizer)	(Page level)	(Optimizer)	(Page level)
Q1	38115	1549	34320	96 %	10 %
Q2	44186	2174	37993	95 %	13 %

(b) All factors, comparison with page level (Countries)

Query Set	Response Time	Response Time	Response Time	%improvement	%improvement
	(No optimization)	(with Optimizer)	(Page level)	(Optimizer)	(Page level)
Q1	22013	1644	17832	93 %	19 %

(c) All factors, comparison with page level (TheaterLoc)

Experimental Results



Total work done with and without materialization
Varied update frequency

Related Work

- Semantic Caching [Dar 1996, Keller and Basu 1996]
 - Compact description is solution to containment problem
- Caching for Web Proxy Servers [Chankunthod et al., 1995]
 - Significant improvement over page level
- Caching for Federated Databases [Goni et al, 1996]
 - Only fixed classes
- View Selection for materialization in Data Warehousing [Gupta and Mumick 1998]
 - Fixed set of views, minimizing cost of data operations
- Extracting patterns from queries similar to mining association rules [Agrawal, Imielinski & Swami 1993]
 - Differences in language learnt, also we have dynamic ontology

Related Work

- Studied applicability of materialization approach in Information Manifold, InfoMaster, TSIMMIS, InfoSleuth, DISCO and Garlic
- Different modeling and query processing approaches
 - views, object-based, datalog based, KR based etc.
- ♦ Materialization framework applicable
- Compact description needed in all
- ◆ Can be extended for source structure analysis
- ◆ Approach for updating data will also apply

Conclusion

Contributions

- Approach for performance optimization by materialization
- Identifying what to materialize
 - Query distribution analysis
 - Source structure analysis
 - The Update issue
- ♦ Future Work
 - Support for multimedia data
 - Tighter integration with query planner
 - Automating collection of statistics about sources
- Extending contributions to other data management areas
 - Query mining
 - Semantic caching in databases
 - Semi-structured data management