
Reformulating CSPs for Scalability with Application to Geospatial Reasoning

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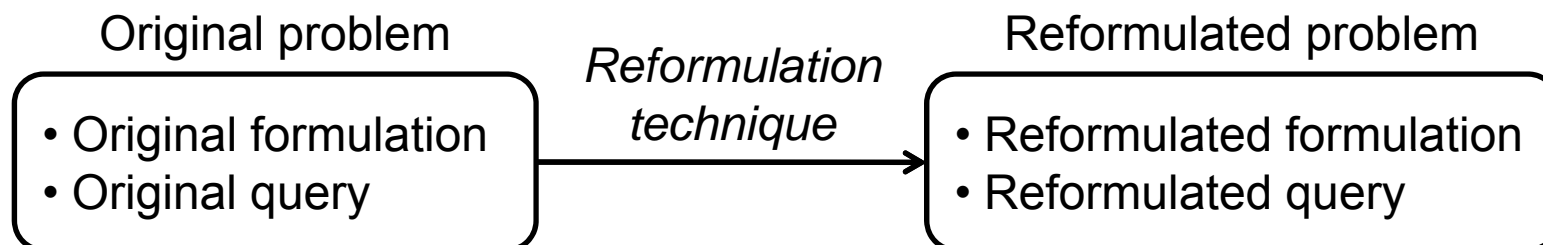
Contributions

- BID problem as a CSP [Michalowski & Knoblock, AAAI 05]
 - Improved constraint model
 - Showed original BID problem is in **P**
 - Custom solver
- Four new reformulation techniques for CSPs
 1. Query reformulation
 2. Domain reformulation
 3. Constraint relaxation
 4. Reformulation via symmetry detection
- Applying the reformulations to the BID problem

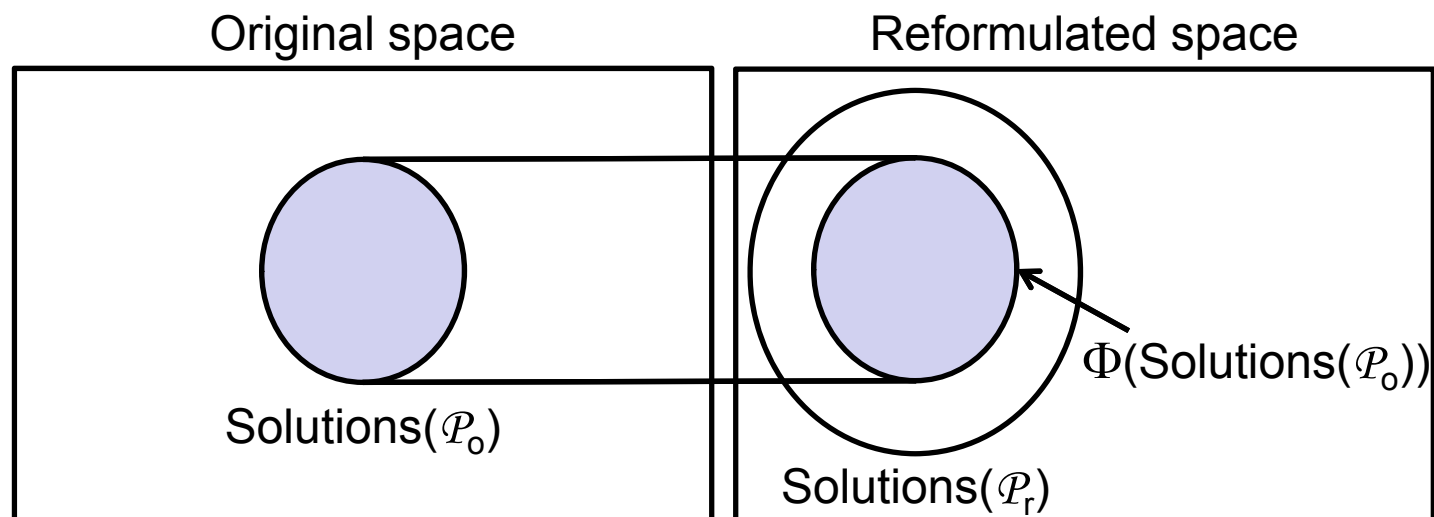
Outline

- **Background**
- BID: CSP model & custom solver
- Reformulation techniques
 - Description
 - General use in CSPs
 - Application to BID
 - Evaluation on real-world BID data
- Conclusions & future work

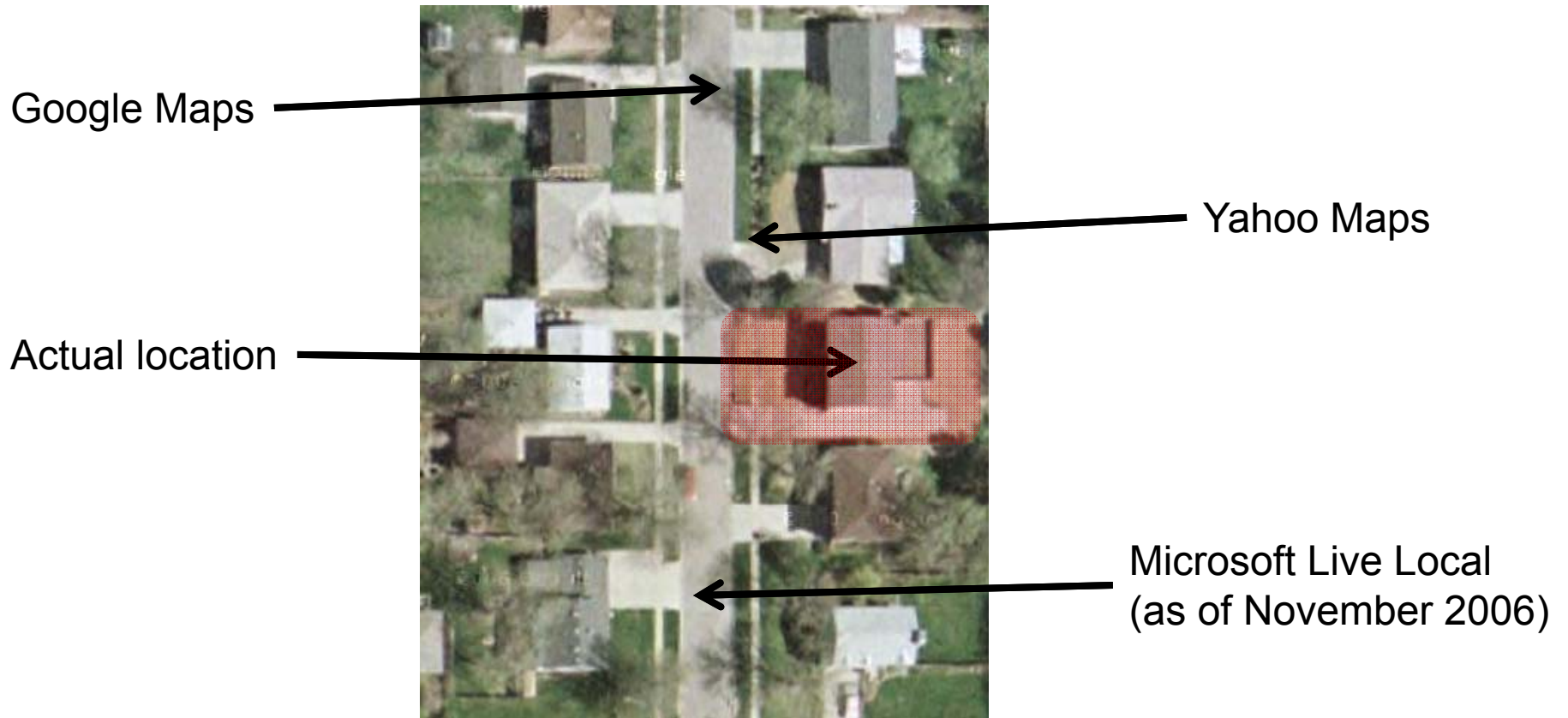
Abstraction & Reformulation



The reformulation may be an approximation

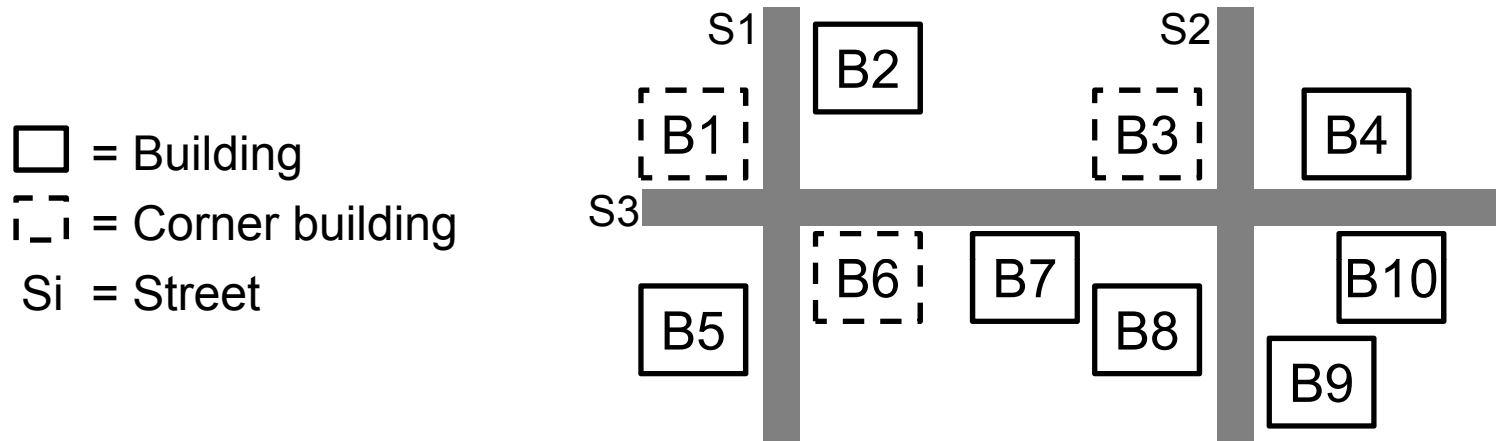


Issue: finding Ken's house



Building Identification (BID) problem

- Layout: streets and buildings

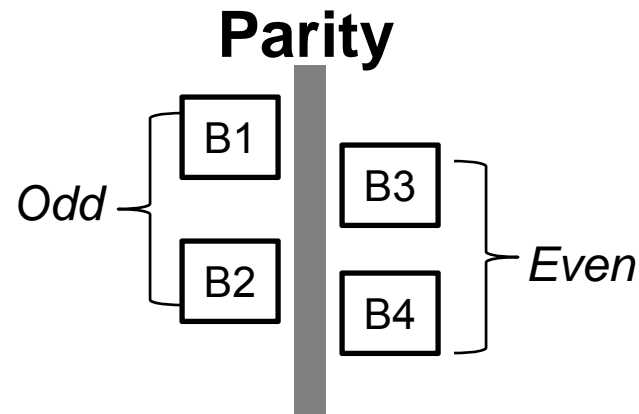
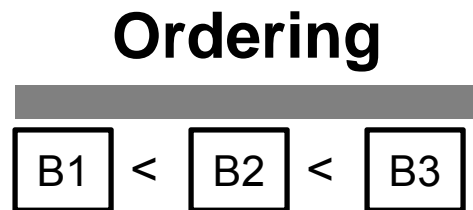


- Phone book
 - Complete/incomplete
 - Assumption: all addresses in phone book correspond to a building in the layout

S1#1, S1#4, S1#8,
S2#7, S2#8, S3#1,
S3#2, S3#3, S3#15,
...

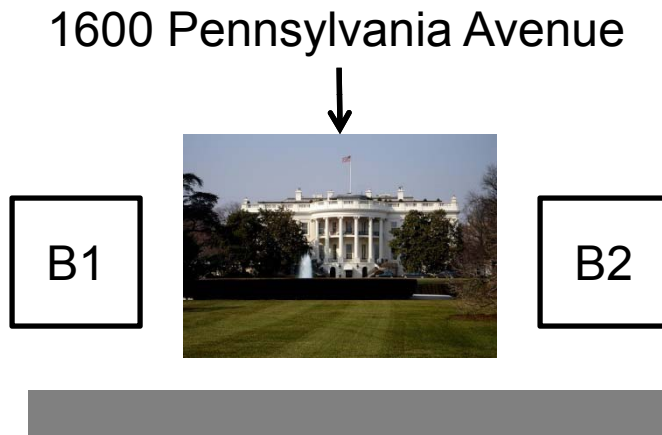
Basic (address numbering) rules

- Ordering
 - Numbers increase/decrease along a street
- Parity
 - Numbers on a given side of a street are odd/even

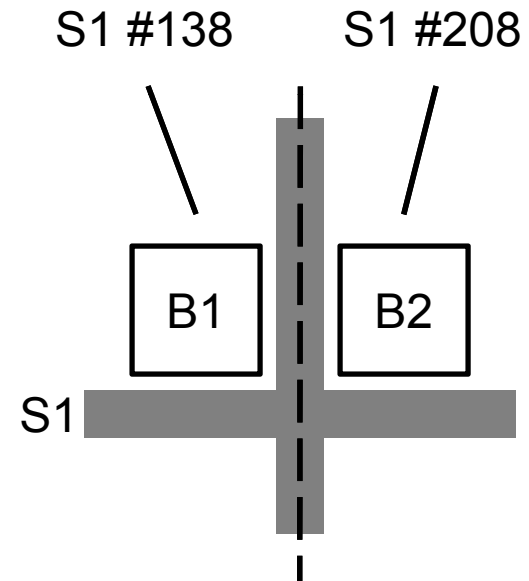


Additional information

Landmarks



Gridlines



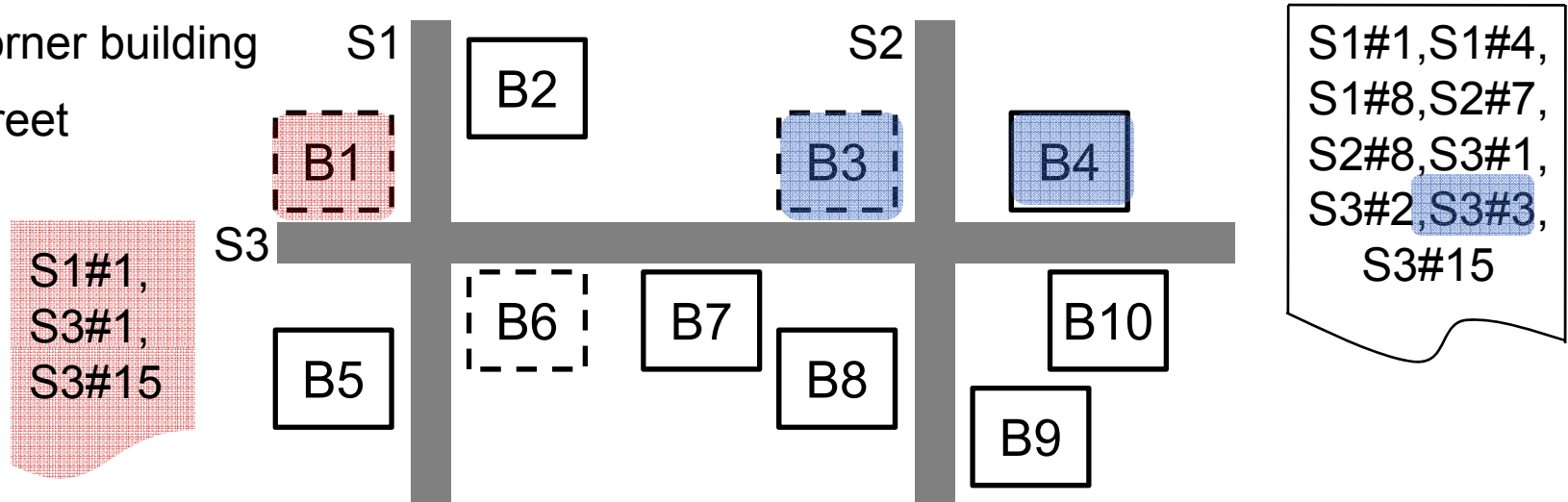
Query

1. Given an address, what buildings could it be?
2. Given a building, what addresses could it have?

□ = Building

□ (dashed) = Corner building

S_i = Street



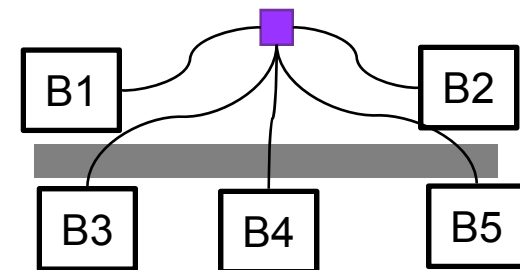
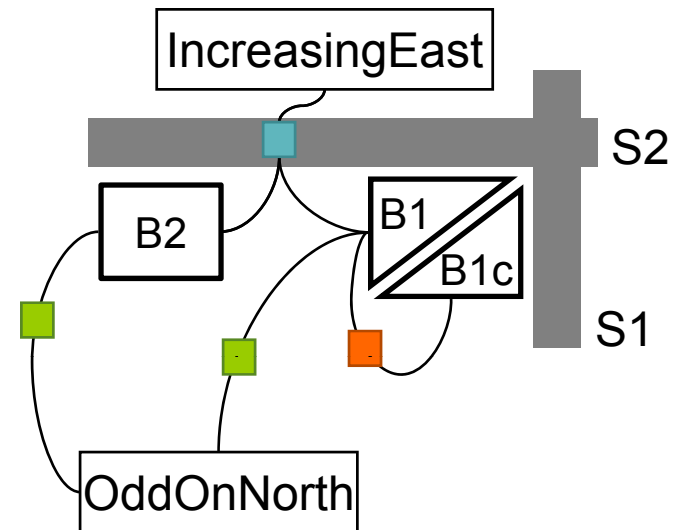
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- **BID model & custom solver**
- Reformulation techniques
- Conclusions & future work

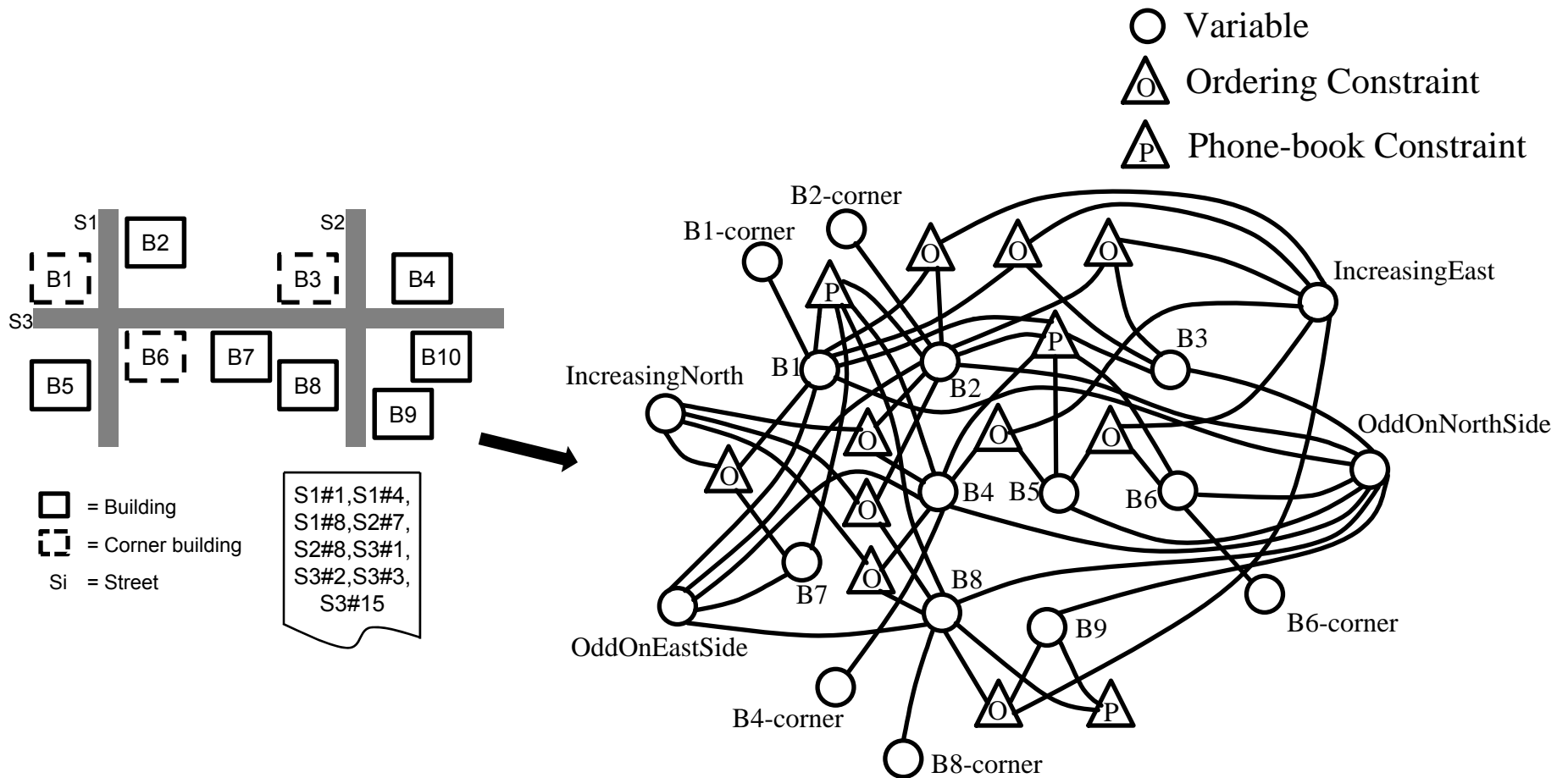
CSP model

- Parity constraints
- Ordering constraints
- Corner constraints

- Phone-book constraints
- Optional: grid constraints



Example constraint network



Features of new model & solver

Improvement over previous work: [Michalowski +, 05]

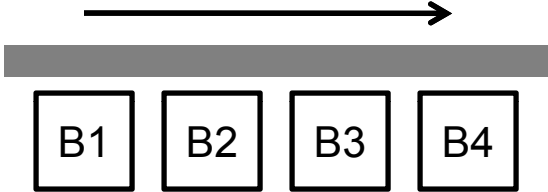
- Model
 - Reduces number of variables and constraints arity
 - Reflects topology: Constraints can be declared locally & in restricted ‘contexts,’ important feature for Michalowski’s work
- Solver
 - Exploits structure of problem (backdoor variables)
 - Implements domains as (possibly infinite) intervals
 - *Incorporates all reformulations (to be introduced)*

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Query in the BID

- Problem: BID instances have many solutions

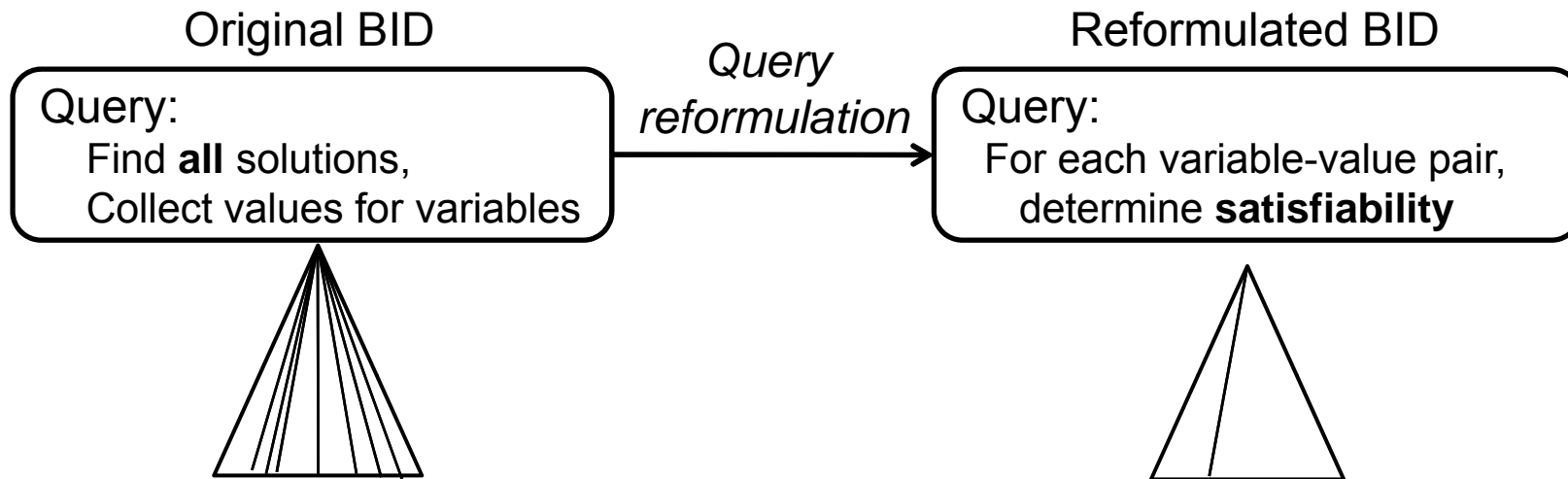


Phone book: {4,8}

B1	B2	B3	B4
2	4	6	8
2	4	8	10
2	4	8	12
4	8	10	12
4	6	8	10
4	6	8	12

We **only** need to know which values (address) appear in **at least one** solution for a variable (building)

Query reformulation



Original query	Reformulated query
Single enumeration problem	Many satisfiability problems
All solutions	Per-variable solution
Exhaustive search	One path
Impractical when there are many solutions	Costly when there are few solutions

Evaluations: real-world data from El Segundo

[Shewale]

Case study	Phone book	Number of...		
	Completeness	Buildings	Corner buildings	Blocks
NSeg125-c	100.0%	125	17	4
NSeg125-i	45.6%			
NSeg206-c	100.0%	206	28	7
NSeg206-l	50.5%			
SSeg131-c	100.0%	131	36	8
SSeg131-i	60.3%			
SSeg178-c	100.0%	178	46	12
SSeg178-i	65.6%			

Previous work did not scale up beyond 34 7 1

Evaluation: query reformulation

Incomplete phone book → many solutions → better performance

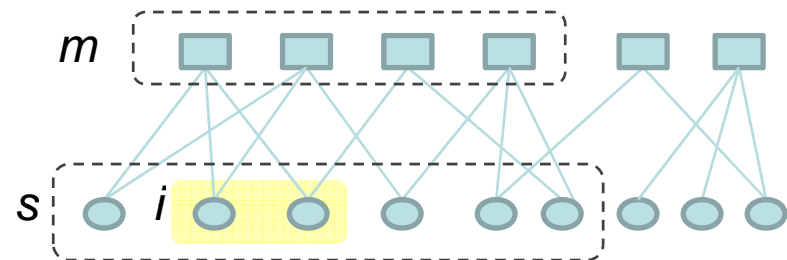
Case study	Original query	New query [s]
NSeg125-i	>1 week	744.7
NSeg206-i	>1 week	14,818.9
SSeg131-i	>1 week	66,901.1
SSeg178-i	>1 week	119,002.4

Complete phone book → few solutions → worse performance

Case study	Original query [s]	New query [s]
NSeg125-c	1.5	139.2
NSeg206-c	20.2	4,971.2
SSeg131-c	1123.4	38,618.4
SSeg178-c	3291.2	117,279.1

Generalizing query reformulation

- Relational (i,m) -consistency, algorithm $R(i,m)C$
 - For every m constraints
 - Compute **all solutions** of length s
 - To generate tuples of length i
 - Space: $O(d^s)$

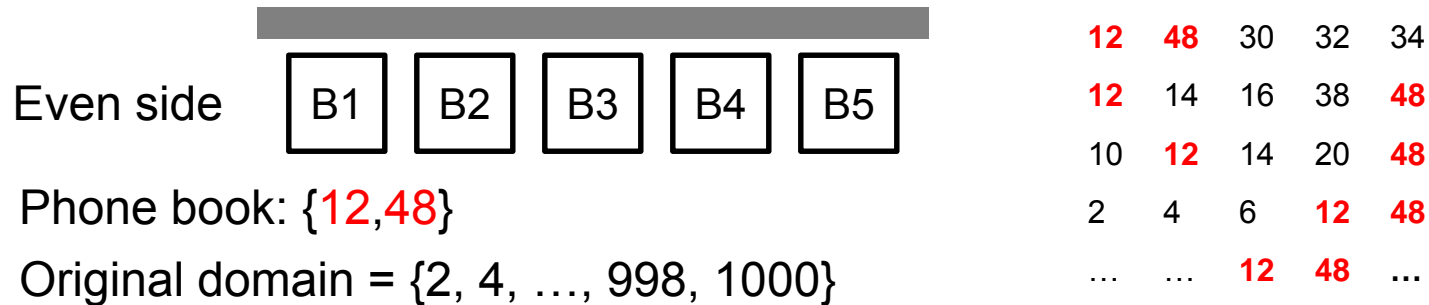


- Query reformulation for Relational (i,m) -consistency
 - For each combination of values for i variables
 - Try to extend to **one** solution of length s
 - Space: $O(\binom{s}{i}d^i)$, $i < s$
- Reformulated BID query is $R(1,|C|)C$

Outline

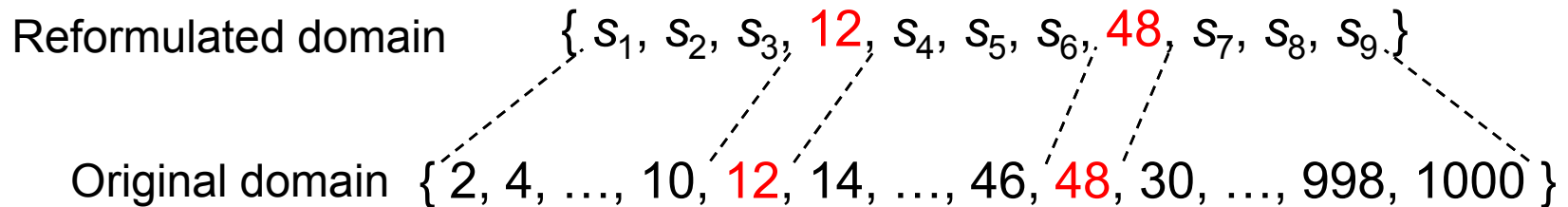
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AllDiff-Atmost in the BID



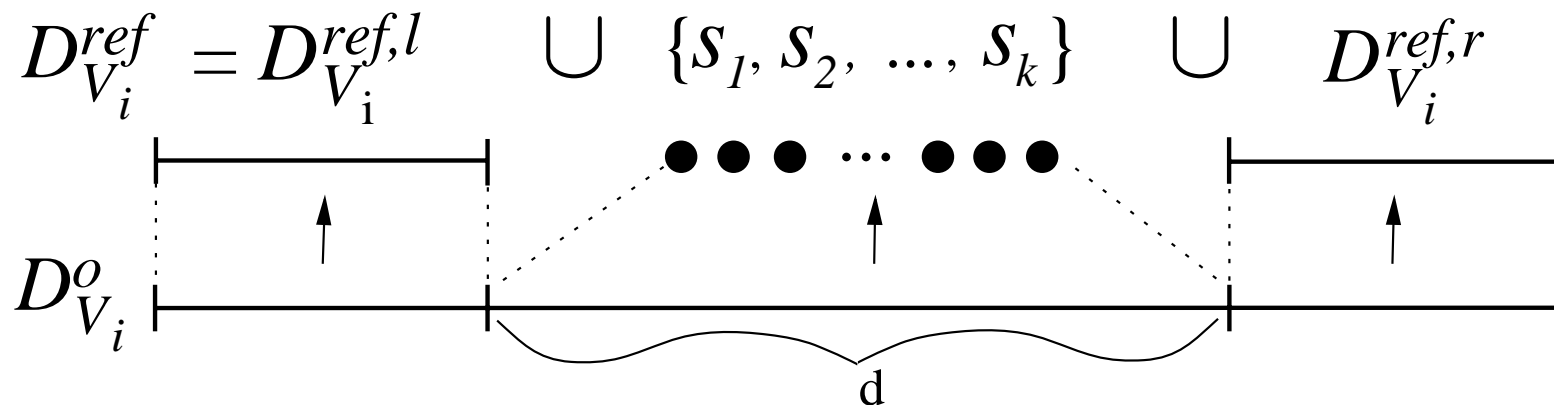
- Can use at most

- 3 addresses in [2,12) AllDiff-Atmost({B1,B2,...,B5},3,[2,12))
- 3 addresses in (12,48) AllDiff-Atmost({B1,B2,...,B5},3,(12,48))
- 3 addresses in (48,1000] AllDiff-Atmost({B1,B2,...,B5},3,(48,1000))



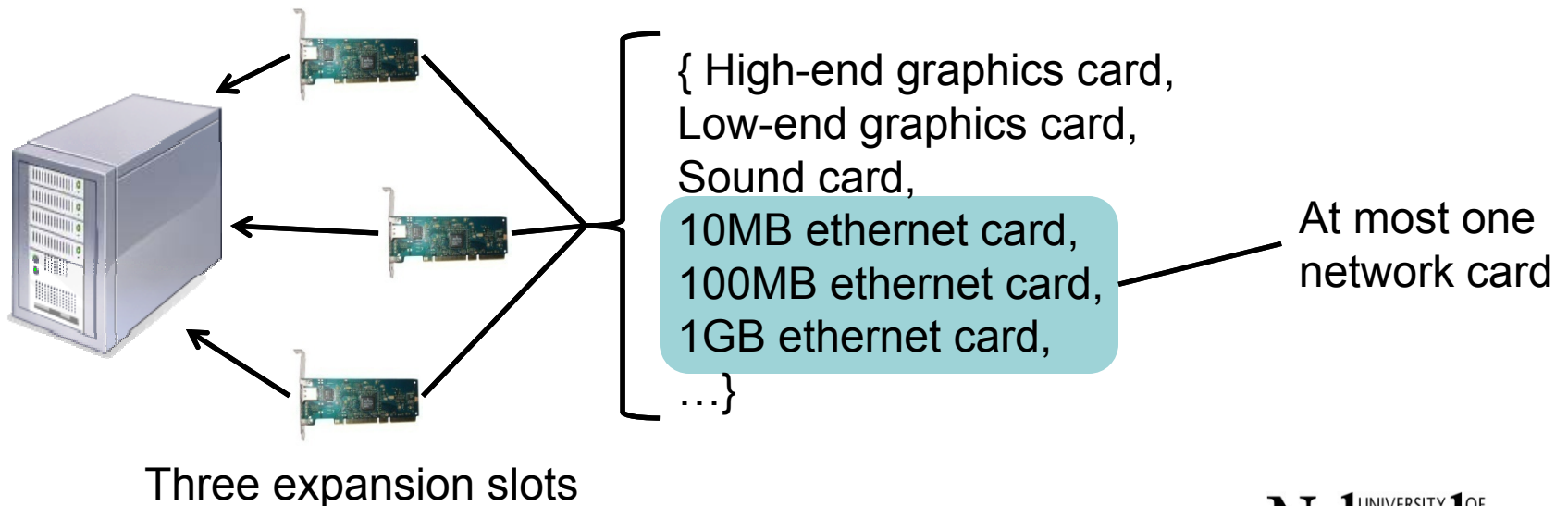
AllDiff-Atmost reformulation

- Given AllDiff-Atmost(\mathcal{A}, k, d)
 - The variables in \mathcal{A} can be assigned at most k values from the set d
- Replace
 - interval d of values (potentially infinite)
 - with k **symbolic values**



AllDiff-Atmost constraint

- AllDiff-Atmost(\mathcal{A}, k, d)
 - The variables in \mathcal{A} can be assigned at most k values from the set d



Evaluation: domain reformulation

- Reduced domain size → improved search performance

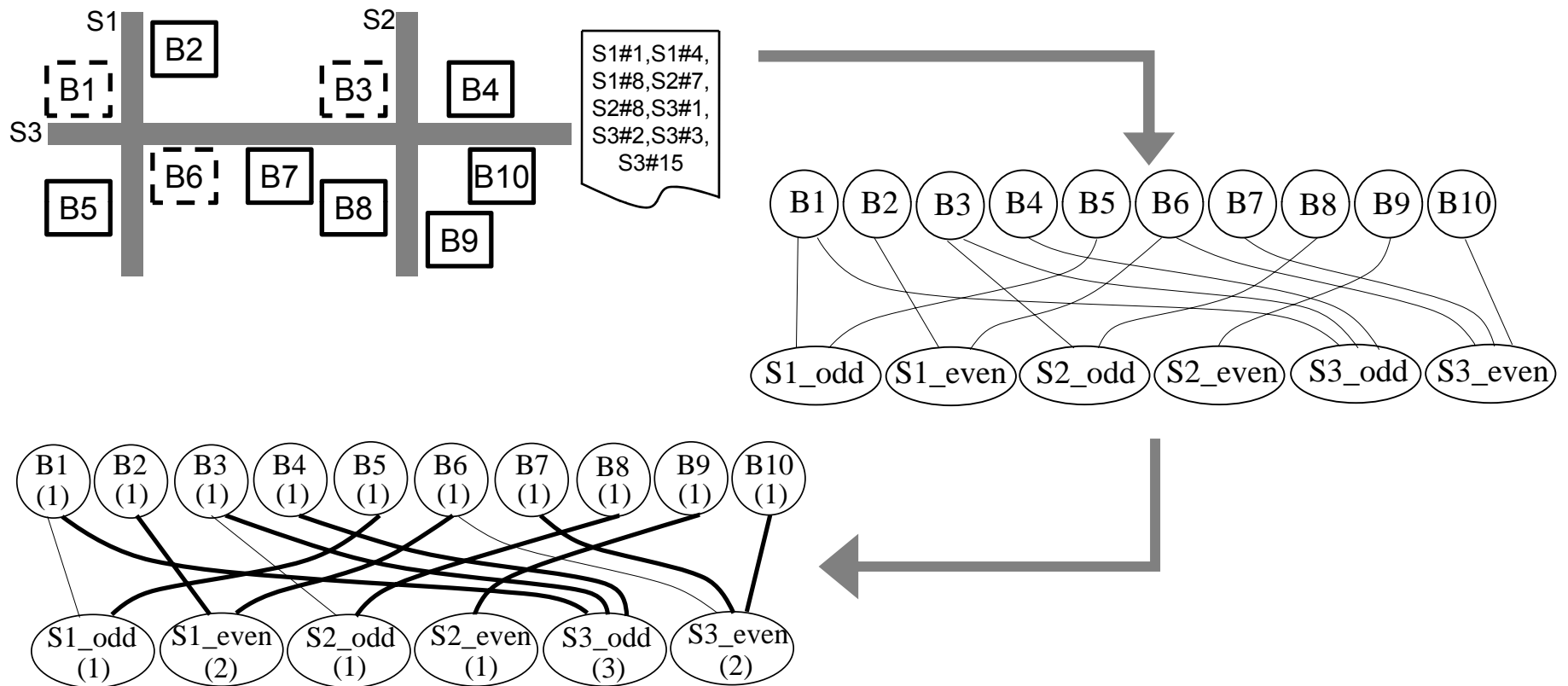
Case study	Phone-book completeness	Average domain size		Runtime [s]	
		Original	Reformulated	Original	Reformulated
NSeg125-i	45.6%	1103.1	236.1	2943.7	744.7
NSeg206-i	50.5%	1102.0	438.8	14,818.9	5533.8
SSeg131-i	60.3%	792.9	192.9	67,910.1	66,901.1
SSeg178-i	65.6%	785.5	186.3	119,002.4	117,826.7

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BID as a matching problem

- Assume we have no grid constraints



BID w/o grid constraints

- BID instances without grid constraints can be solved in *polynomial time*

Case study	Runtime [s]	
	BT search	Matching
NSeg125-c	139.2	4.8
NSeg206-c	4971.2	16.3
SSeg131-c	38618.3	7.3
SSeg178-c	117279.1	22.5
NSeg125-i	744.7	2.5
NSeg206-i	5533.8	8.5
SSeg131-i	38618.3	7.3
SSeg178-i	117826.7	4.9

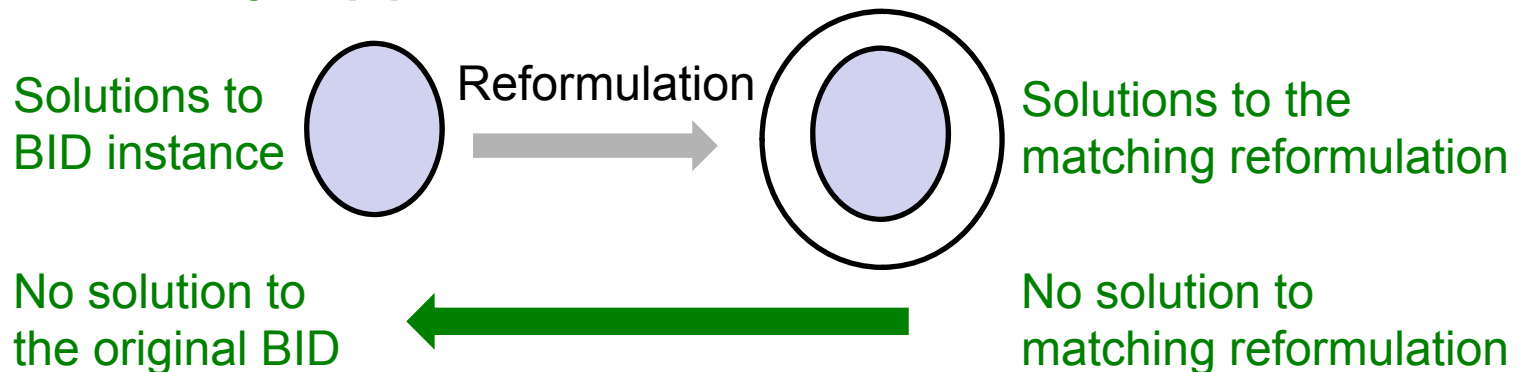
BID w/ grid constraints

1. Filter CSP

[Régis, 1994]

Remove vyps that do not appear in a maximum matching

2. Relaxed CSP: matching reformulation as a *necessary approximation* of the BID



Matching reformulation in Solver

Filter CSP..

Preproc1

For every vvp

Consider CSP + vvp

If relaxed CSP is solvable

Preproc2

Find one solution using BT search

At each instantiation, filter CSP

Lookahead

Evaluation: matching reformulation

- Generally, improves performance

Case Study	BT	Preproc2 +BT	% (from BT)	Lkhd +BT	% (from BT)	Lkhd +Preproc1&2 + BT	% (from Lkhd+BT)
NSeg125-i	1232.5	1159.1	6.0%	726.6	41.0%	701.1	3.5%
NSeg206-c	2277.5	614.2	73.0%	1559.2	31.5%	443.8	71.5%
SSeg178-i	138404.2	103244.7	25.4%	121492.4	12.2%	85185.9	29.9%

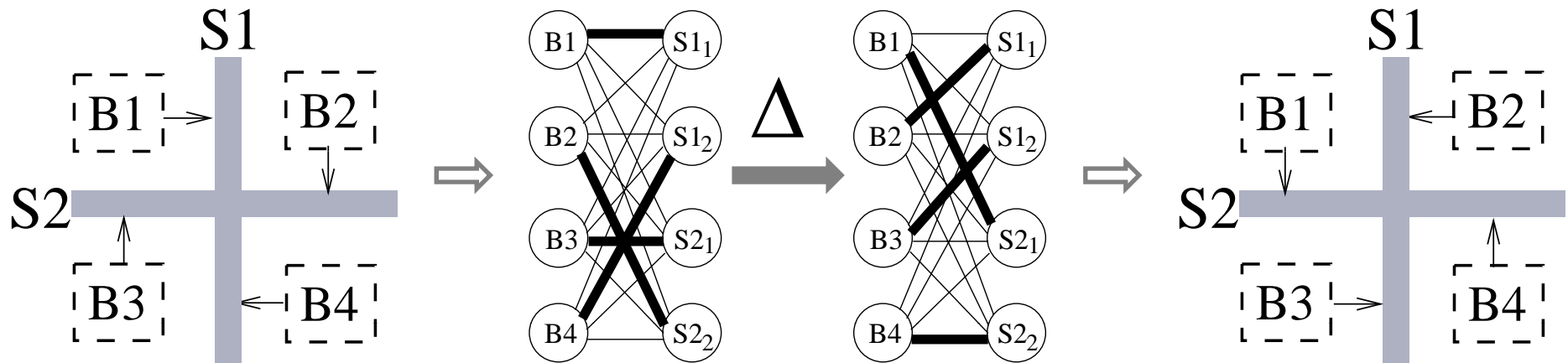
- Rarely, the overhead exceeds the gains

Case Study	BT	Preproc2 +BT	% (from BT)	Lkhd +BT	% (from BT)	Lkhd +Preproc1&2 + BT	% (from Lkhd+BT)
NSeg125-c	100.8	33.2	67.1%	140.2	-39.0%	29.8	78.7%
NSeg131-i	114405.9	114141.3	0.2%	107896.3	5.7%	108646.6	-0.7%

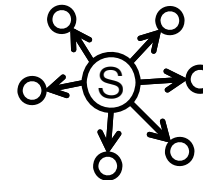
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Symmetric matchings in BID



- **All** matchings can be produced from the symmetric difference of
 - a single matching and
 - a set of disjoint alternating cycles & paths starting @free vertex
- Some symmetric solutions do not break grid constraints
 - Ignore symmetric solutions during search
- Some do, we do not know how to use them...



Conclusions

- We proposed four reformulation techniques
- We described their usefulness for general CSPs
- We demonstrated their effectiveness on the BID

Lesson:

Reformulation is an effective approach to improve the scalability of complex combinatorial systems

Future work

- Empirically evaluate our new algorithm for relational (i,m) -consistency
- Exploit the symmetries we identified
- Enhance the model by incorporating new constraints [Michalowski]

Questions?